



### **Are Some Voters More Equal Than Others?**

#### **Discussions and work in progress on formalisation**

Hugo Jonker

in collaboration with Ben Smyth



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Elections should be fair.



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#### computation: every path must occur in an infinite computation.

#### contract signing:

Either all or none of the parties receive a signed document.

#### two-party exchange:

Either both items change owner, or neither does.

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Each voter has specific, partial control over the result. Fairness is broken when a voter can exercise control beyond this. Control: "+1"? can vary per voter?



### **Examples for discussion**

## Discussion: ballot independence

#### Situation 1 (copy ballot):

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## Discussion: ballot independence

#### Situation 1 (copy ballot):

#### Submit a copy of another voter's filled in ballot.

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#### Situation 2 (vote unlike someone):

Submit a modified copy of another voter's filled in ballot.

- You can vote unlike someone.
- Privacy problem? Fairness problem?

## Discussion: aborting a vote

Situation 3 (*change your mind*):

If voting occurs in > 1 phase, don't participate in last phase.

- You can cancel your vote.
- When is this a fairness problem?



### **Existing formalisations**

[KR05]: to verify [FOO92].

- no one can learn vote v before opening phase.
   Standard ProVerif secrecy check of vote variable v.
- no one can guess v before opening phase.  $\phi \approx_s \nu v.\phi$  ProVerif check.

- + automatic checking
- copying/modifying ballot not caught
- contents of vote  $\stackrel{?}{\Longrightarrow}$  no early results



[BRS07]: to verify [FOO92].

$$\neg$$
resultAnnounced  $\implies \bigwedge_{a \in Ag} L_a(\bigwedge_{b \neq a, c \in \mathcal{C}} vote_b(c)).$ 

Before results, no one can exclude any choice by any other voter.

- + knowledge based reasoning
- + straightforward definition
- how to apply
- fairness > knowing no ballots

[TMT<sup>+</sup>08]: case study of [FOO92].

$$\nu X. \wedge_{c \in \mathcal{C}} \left( \langle x. (x_i - X_s \triangleright X_r : v). y \leftrightarrow \varepsilon \rangle tt \rightarrow \\ \langle x. d(T). y. (x_i - X_s \triangleright X_r : v). z \leftrightarrow x. d(T). y. z \rangle X \right) \right)$$

If a vote can be determined, then there must have been a phase boundary earlier in the protocol.

- "normalized" protocol
- non-intuitive language
- guessing attacks not caught
- ballot exposure  $\neq$  fairness

## BHM08: don't re-use the vote

[BHM08]: general def of "soundness", applied to [JCJ05].

Every eligible voter votes once.

- $t = t1 \cdot \operatorname{start}(id) \cdot t2$ .
- **Eligibility:** start(*id*)  $\notin t1 \cdot t2$ .
- One vote:  $newid(id) \in t1$ . (event by id manager).

- + simple, straightforward def
- limited to "soundness" / accuracy + democracy



### **Towards formalising fairness**



#### If the result is unaffected, fairness is not harmed.

- 1. before voting, voter observes trace t;
- 2. t can be extrapolated to full run with and without voter;
- 3. For all such possible extrapolations: determine result;

Fairness:  $\exists c \colon \forall t \in Tr(with) \colon \exists t' \in Tr(without) \colon c =$ result(t) -result $(t') \land t \approx t'$ .



- Constant = 1,2,3,...
- Difference between two results constant necessary? Sufficient?
- What if no one votes after voter? Or a variable number?



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- b.ii. For every voter, the effect should be the same: A change of 1 vote.
  - c. The voting system does not influence the vote.
  - d. No pulling out (problem in FOO).

# Conclusions

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Thank you for your attention. Questions/comments?

# **References**

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